

Last time:

- optimal control
- dynamic programming
(discrete-time)

Lecture 4
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This time:

- dyn. prog. (cont. time)
- multi-agent games
- robust opt. ctrl.

Dynamic Programming - Continuous Time:

One of the advantages of D.P. is that it can be used to solve both discrete & continuous-time optimal control problems.

In continuous-time, principle of optimality enables:

$$V(x, t) := \min_{u(\cdot) \in \mathcal{U}_t^T} \left\{ \int_{\tau=t}^T L(x(\tau), u(\tau)) d\tau + l(x(T)) \right\}$$

"all control signals from $[t, T]$ "

$$= \min_{u(\cdot) \in \mathcal{U}_t^T} \left\{ \int_{\tau=t}^{t+\delta} L(x(\tau), u(\tau)) d\tau + \int_{s=t+\delta}^T L(x(s), u(s)) ds + l(x(T)) \right\}$$

$$= \min_{u(\cdot) \in \mathcal{U}_t^{t+\delta}} \left\{ \int_{\tau=t}^{t+\delta} L(x(\tau), u(\tau)) d\tau \right\} + \min_{u(\cdot) \in \mathcal{U}_{t+\delta}^T} \left\{ \int_{s=t+\delta}^T L(x(s), u(s)) ds + l(x(T)) \right\}$$

cost from $[t, t+\delta]$:= $V(x(t+\delta), t+\delta)$
By def. of optimal cost from $t+\delta$ onwards

$$= \min_{u(\cdot) \in \mathcal{U}_t^{t+\delta}} \left\{ \int_{\tau=t}^{t+\delta} L(x(\tau), u(\tau)) d\tau + V(x(t+\delta), t+\delta) \right\}$$

lets focus on studying what happens when $\delta \rightarrow 0$. For now, assume $V(\cdot, \cdot)$ is everywhere differentiable. lets go through an informal derivation based on finite-element analysis of how V changes when $\delta > 0$ but is very small.

$$V(x, t) = \min_{u(\cdot) \in \mathcal{U}_t^{t+\delta}} \left\{ \int_{\tau=t}^{t+\delta} L(x(\tau), u(\tau)) d\tau + V(x(t+\delta), t+\delta) \right\}$$

$$V(x, t) \approx \min_{u(t) \in \mathcal{U}} \left\{ L(x(t), u(t)) \cdot \delta + \underbrace{V(x(t+\delta), t+\delta)} \right\}$$

lets simplify by taking Taylor Series Expansion of $V(x(t+\delta), t+\delta)$ around current $(x(t), t)$.

$$V(x(t+\delta), t+\delta) \approx V(x(t), t) + \frac{\partial V}{\partial x} \cdot \underbrace{(x(t+\delta) - x(t))}_{= f(x(t), u(t)) \cdot \delta} + \frac{\partial V}{\partial t} \cdot (t+\delta - t) + \text{h.o.t.}$$

$$\approx f(x(t), u(t)) \cdot \delta$$

Plugging expansion of $V(x(t+\delta), t+\delta)$ back in:

$$V(x, t) \approx \min_{u(t) \in U} \left\{ L(x(t), u(t)) \cdot \delta + \underbrace{V(x(t), t)}_{\text{boxed}} + \frac{\partial V}{\partial x} \cdot \underbrace{f(x(t), u(t)) \cdot \delta}_{\text{boxed}} + \frac{\partial V}{\partial t} \cdot \underbrace{(t+\delta - t)}_{\text{boxed}} + \text{h.o.t.} \right\}$$

$$\cancel{V(x, t)} = \cancel{V(x(t), t)} + \frac{\partial V}{\partial t} \cdot \delta + \min_{u(t) \in U} \left\{ L(x(t), u(t)) \cdot \delta + \frac{\partial V}{\partial x} \cdot f(x(t), u(t)) \cdot \delta \right\}$$

$$0 = \frac{\partial V}{\partial t} \cdot \delta + \min_{u(t) \in U} \left\{ L(x(t), u(t)) \cdot \delta + \frac{\partial V}{\partial x} \cdot f(x(t), u(t)) \cdot \delta \right\}$$

$$0 = \delta \left[\frac{\partial V}{\partial t} + \min_{u(t) \in U} \left\{ L(x(t), u(t)) + \frac{\partial V}{\partial x} \cdot f(x, u) \right\} \right]$$

↪ must hold true for all $\delta > 0$, can divide both sides by δ to get:

Hamilton-Jacobi-Bellman (HJB) Equation

$$\frac{\partial V}{\partial t} + \min_{u(t) \in U} \left\{ L(x(t), u(t)) + \frac{\partial V}{\partial x} \cdot f(x, u) \right\} = 0$$

$V(x, T) = l(x)$ called Hamiltonian

⊛ continuous-time counterpart of Bellman Equation

⊛ terminal time PDE

⊛ optimal control:

$$u^*(x, t) = \arg \min_{u \in U} \left\{ L(x, u) + \frac{\partial V}{\partial x} \cdot f(x, u) \right\}$$

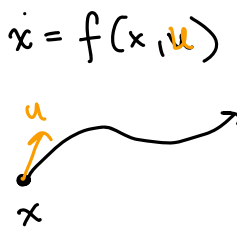
! key equations that we will solve during safety analysis!

let's add uncertainty back in! For now, we will focus on non-deterministic uncertainty, $d \in \mathcal{D}$, but I want to note that in the paper discussions + later in the course, we will see probabilistic uncertainty.

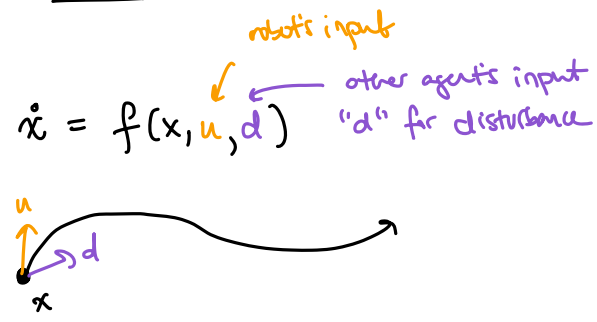
we will formulate this as a robust problem. More specifically, as a

Zero-sum dynamic game
 there is winner & a loser game evolves over time result of the interaction depends on 2+ players

Recall, before:



Now:



Maybe, one way we could pose the optimization problem:

$$V(x, t) = \min_{u(\cdot) \in \mathcal{U}_t^T} \max_{d(\cdot) \in \mathcal{D}_t^T} J(x, u(\cdot), d(\cdot), t)$$

s.t. $\dot{x}(\tau) = f(x(\tau), u(\tau), d(\tau))$, $\forall \tau \in [t, T]$
 $u(\tau) \in \mathcal{U} = \mathbb{R}^m$
 $d(\tau) \in \mathcal{D} = \mathbb{R}^d$

INFORMATION PATTERNS:

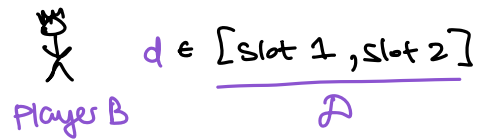
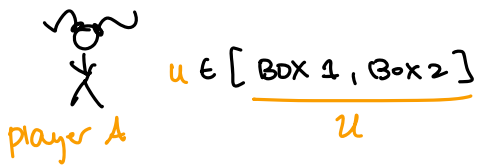
When we have 2 players reacting to each other, their optimal strategy will depend on what information they each have access to.

EXAMPLE Suppose there are 2 boxes, each with 2 slots.

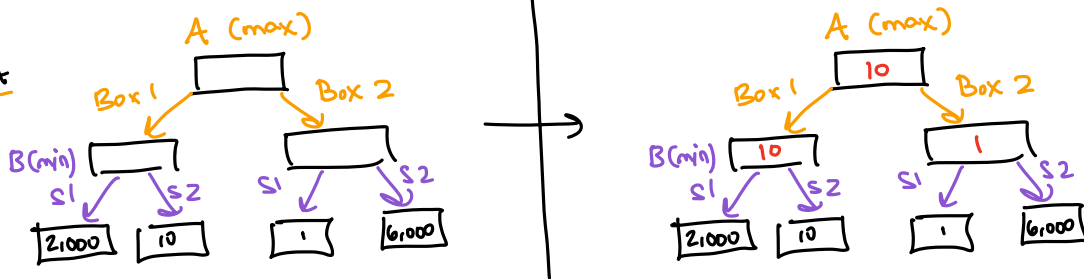
Each slot contains some prize money. Player A (you :) wants to maximize the prize money, while Player B (e.g. competition org.) wants to minimize Player A's prize money.

Box 1	
s1	s2
2,000	10

Box 2	
s1	s2
1	6,000



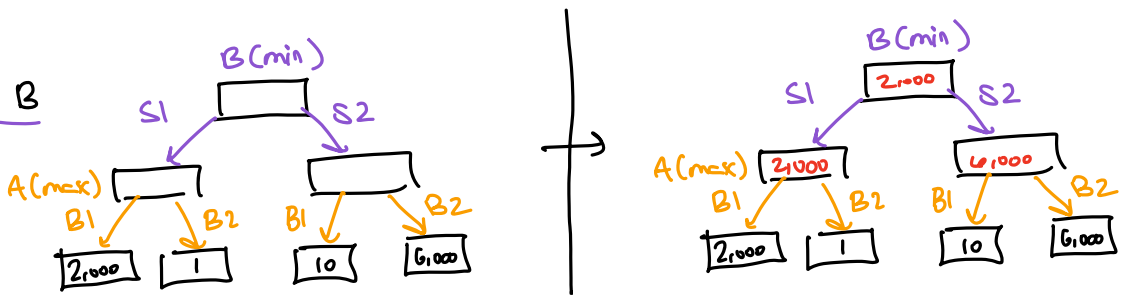
Suppose Player A goes first:



⇒ Best outcome for A is to pick Box 1, and get \$10.

$$\max_{u \in [B1, B2]} \left(\min_{d \in [s1, s2]} J(u, d) \right) = 10 \quad \leftarrow \text{Mathematically}$$

Suppose Player B goes first:



Best strategy of Player B is to choose slot 1 and pay player A a reward of \$2,000.

$$\min_{d \in [s1, s2]} \left(\max_{u \in [B1, B2]} J(u, d) \right) = 2,000$$

This can be stated formally via the minimax inequality:

$$\max_a \left(\min_b J(a, b) \right) \leq \min_b \left(\max_a J(a, b) \right)$$

just talked about order of play what about this?

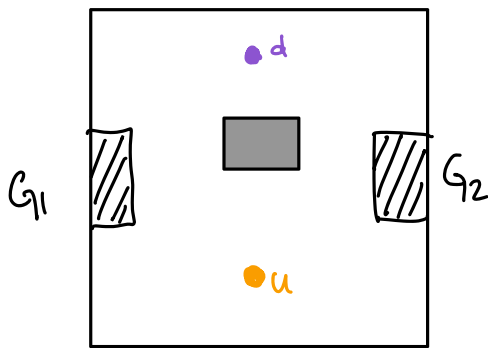
In dynamic games, the outcome depends on WHEN and WITH WHAT INFO.
 Each player decides their inputs.

In fact, earlier I wrote one kind of information pattern:

OPEN-LOOP INFORMATION PATTERN

$$V(x|t) = \min_{u(\cdot) \in \mathcal{U}_t^T} \max_{d(\cdot) \in \mathcal{D}_t^T} J(x, u(\cdot), d(\cdot), t)$$

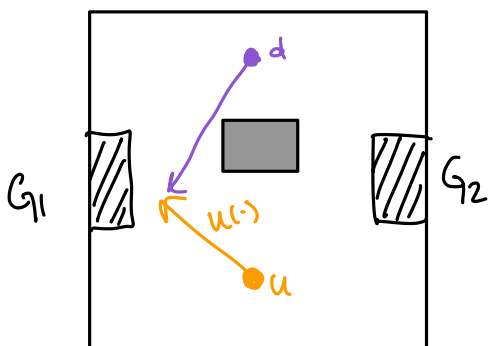
Let's consider example.



Here, u wants to get to G1 or G2, without interception from d . But d wants to intercept u .

1) OPEN-LOOP INFO. PATTERN

$$V(x|t) = \min_{u(\cdot) \in \mathcal{U}_t^T} \max_{d(\cdot) \in \mathcal{D}_t^T} J(x, u(\cdot), d(\cdot), t)$$



$u(\cdot)$ will declare its entire control signal $[t, T]$

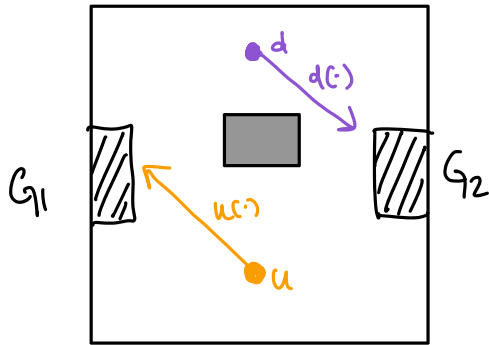
$d(\cdot)$ gets to see this entire signal before choosing its own.

" d can see "future" and u cannot change their mind"

⇒ OVERLY PESSIMISTIC

SWAP ORDER OF PLAY:

$$V(x,t) = \max_{d(\cdot) \in \mathcal{D}_t^T} \min_{u(\cdot) \in \mathcal{U}_t^T} J(x, u(\cdot), d(\cdot), t)$$



$d(\cdot)$ will commit / declare its entire control signal from $[t, T]$

$u(\cdot)$ gets to see this before choosing its own.

⇒ overly optimistic

2) CLOSED-LOOP (FEEDBACK) INFORMATION PATTERNS

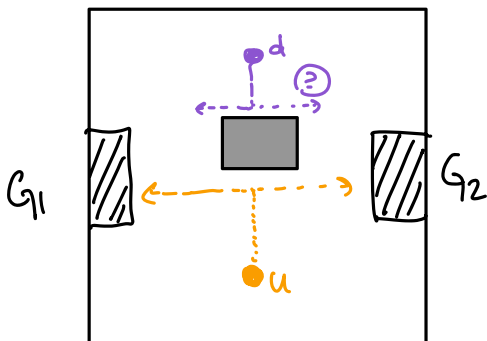
These above formulations are not suitable for most practical systems.

We would like the controller to adapt as the interaction evolves, BUT respect the fact that @ time t , we only have information up to time t .

In differential games, such strategies are called

NON-ANTICIPATIVE STRATEGIES:

$$V(x,t) = \min_{u(\cdot) \in \Gamma_u^T} \max_{d(\cdot) \in \Gamma_d^T} J(x, u(\cdot), d(\cdot), t)$$



In this case, $u(\cdot)$ does not have to declare that they want to go left or right @ start of interaction.

Here, $d(\cdot)$ cannot simply intercept $u(\cdot)$ before it reaches goal

Definition: Non-anticipative strategies:

$$\Gamma_d(t, T) = \left\{ \delta : \mathcal{U}_t^T \rightarrow \mathcal{D}_t^T \text{ s.t. if } \forall u_1(\cdot), u_2(\cdot) \in \mathcal{U}_t^T, \forall \tau \in [t, T] \right.$$

$$\left. \begin{aligned} & (u_1(s) = u_2(s) \text{ a.e. } s \in [t, \tau]) \Rightarrow \\ & \left(\delta[u_1](s) = \delta[u_2](s) \text{ a.e. } s \in [t, \tau] \right) \end{aligned} \right\}$$

"if the controls are the same almost everywhere..."
 "then the disturbance must react the same way over that t-hor."

- δ is a map from u's control signals to d's ctrl. signals
- Intuition: the disturbance δ cannot pre-emptively start to adapt to a change in u UNTIL SUCH A CHANGE BEGINS!

